

# The relative approximation degree I

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Where are we?

we want to prove (OIMM) / Henselian rationality

## Proof structure

- ① rank 1 + separably closed ← we are still here
  - ② rank 1
  - ③ finite rank
  - ④ arbitrary rank
- Fouzei outlined the proof two weeks ago

black boxes:

- Galois-degree-p-extensions → Lemma 4.8 (Artin-Schreier-extension / equi-p, Sylvy last week)
- Lemma 4.9 (Kummer extensions / mixed char, ???)

• **TODAY:** Kuhlmann-Vlaha, Theorem 11.1  
(+ prob. more)

$(K, v)$  valued field, rank 1,

$(F|K, v)$  immediate function field,  $\text{fd} \deg(F|K) = 1$

Suppose there is  $x \in F^h \setminus K^c$  with transcendental approximation type over  $K$  such that  $F^h = K(x)^h$

Then there is already some  $y \in F$  such that  $F^h = K(y)^h$

## § Approximation Types

### Definition

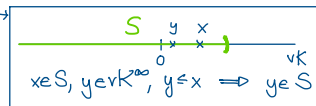
- $B_\alpha(c, K) = \{a \in K : v(a-c) \geq \alpha\}$  "closed" ball in  $(K, v)$  of radius  $\alpha \in vK^\infty = vKv^{-\infty}$

- An approximation type over  $(K, v)$  is a collection

$$\underline{A} = \{B_\alpha(c_\alpha, K) : \alpha \in S\}$$

where  $S \subseteq vK^\infty$  is an initial segment

and the balls  $B_\alpha(c_\alpha, K)$  are linearly ordered by inclusion



- $S$  is called support of  $\underline{A}$ , i.e. no two balls are disjoint!

write  $S = \text{supp } \underline{A}$

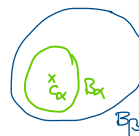
- For  $\alpha \in vK^\infty$ ,  $\underline{A}_\alpha := \begin{cases} B_\alpha(c_\alpha, K) & \text{if } \alpha \in \text{supp } \underline{A} \\ \emptyset & \text{otherwise} \end{cases}$

### Remark

$\underline{A}$  is determined by  $(\underline{A}_\alpha)_{\alpha \in T}$  where  $T \subseteq_{\text{initial}} \text{supp } \underline{A}$

(Because for  $\beta < \alpha \in \text{supp } \underline{A}$ ,

$$\underline{A}_\beta = B_\beta(c_\beta, K) = B_\beta(c_\alpha, K)$$



### Definition

$(L|K, v)$ ,  $x \in L$ . Define

$$\text{appr}(x, K)_\alpha := \{c \in K : v(x-c) \geq \alpha\} = B_\alpha(x, L) \cap K$$

And  $\text{appr}(x, K) := \{\text{appr}(x, K)_\alpha : \alpha \in vK^\infty, \text{appr}(x, K)_\alpha \neq \emptyset\}$ ,

the approximation type of  $x$  over  $K$ .

## 5 Immediate Approximation Types

### Definition

A approximation type over  $(K, v)$ .

•  $(LK, v)$ ,  $x \in L$  Say  $x$  realizes  $\underline{A}$  (in  $(L, v)$ )  $\Leftrightarrow \underline{A} = \text{appr}(x, K)$

•  $\underline{A}$  is trivial  $\Leftrightarrow \underline{A}$  is realized by some  $c \in K$

$$\Leftrightarrow \underline{A}_\infty \neq \emptyset \Leftrightarrow \underline{A}_\infty = \{c\} \Leftrightarrow \text{supp } \underline{A} = vK^\infty$$

•  $\underline{A}$  is immediate  $\Leftrightarrow \bigcap_{\alpha \in \text{supp } \underline{A}} \underline{A}_\alpha = \emptyset$

### Remark (immediate $\Rightarrow$ non-trivial)

$\underline{A}$  trivial  $\Rightarrow \bigcap \underline{A} = \underline{A}_\infty \neq \emptyset \Rightarrow \underline{A}$  not immediate

### Fact/Kaplansky [KV, Proposition 6.6]

Every immediate approximation type is realized in some immediate simple valued field extension

### Lemma [KV, Lemma 4.1]

$(LK, v)$ ,  $x \in L$

(a)  $\text{appr}(x, K)$  is immediate  $\Leftrightarrow v(x-K)$  has no maximal element

(c)  $\text{appr}(x, K)$  is immediate  $\Rightarrow \text{supp } \text{appr}(x, K) = v(x-K)$

Proof: (a)  $\Rightarrow$ :  $\underline{A} = \text{appr}(x, K)$  immediate:  $\bigcap_{\alpha \in \text{supp } \underline{A}} \underline{A}_\alpha = \emptyset$ ,  $c \in K$  arbitrary

$\Rightarrow \exists x \in \text{supp } \underline{A}$ :  $c \notin \underline{A}_\alpha = \mathcal{B}_\alpha(x, L) \cap K \Rightarrow v(x-c) > \alpha$

Let  $c' \in \underline{A}_\alpha = \mathcal{B}_\alpha(x, L) \cap K \neq \emptyset \Rightarrow v(x-c') \geq \alpha > v(x-c)$

So  $v(x-K)$  has no maximal element

$\Leftarrow$ :  $\text{appr}(x, K)$  imm.  $\Rightarrow x \in L \setminus K$  (non-trivial)

Suppose  $v(x-K)$  has no max:  $\forall c \in K \exists c' \in K: v(x-c') > v(x-c)$

To show: for every  $c \in K$  ex.  $\alpha \in \text{supp } \text{appr}(x, K)$ , s.th.  $c \notin \text{appr}(x, K)_\alpha$

Fix  $c \in K$ . There is  $c', c''$  s.th.  $v(x-c'') > v(x-c') > v(x-c)$

$$\stackrel{\Delta\text{-ineq}}{\Rightarrow} v(c'-c) = v((x-c)-(x-c')) = v(x-c) < v(x-c') = v(c''-c')$$

$\Rightarrow v(c''-c) \in v(x-K)$ ,  $c \notin \text{appr}(x, K)_{v(c''-c)} \neq c'$

(c)  $\Leftarrow$ :  $\alpha \in \text{supp } \text{appr}(x, K) \Rightarrow \text{appr}(x, K)_\alpha \neq \emptyset \rightarrow \exists c \in K: v(x-c) \geq \alpha$

-  $v(x-c) = \alpha \Rightarrow \alpha \in v(x-K)$

-  $v(x-c) > \alpha$ : let  $d \in K$  with  $v(d) = \alpha$ , then

$$v(x-(c+d)) = v(d) = \alpha \in v(x-K)$$

$\Leftarrow$ :  $c \in K$ . Since  $v(x-K)$  has no maximal element, there is  $c' \in K$  s.th.

$$v(x-c') > v(x-c)$$

$$\Rightarrow v(c'-c) = v((x-c)-(x-c')) = v(x-c) \in vK$$

$v(x-c) \geq v(x-c) \Rightarrow c \in \text{appr}(x, K)_{v(x-c)} \neq \emptyset \Rightarrow v(x-c) \in \text{supp } \text{appr}(x, K)$

## § Polynomials

### Definition

$\varphi(X)$  formula.  $\Delta$  approximation type,  $(L|K, v)$ ,  $x \in L$ ,  
 $\varphi(X)$  term in  $vK$ .

- write  $\varphi(c)$  for  $c \nearrow \Delta$  if there is  $\alpha \in \text{supp } \Delta$  s.th.  $\varphi(c)$  holds for all  $c \in \Delta_\alpha$
- write " $c \nearrow x$ " for " $c \nearrow \Delta$ " if  $\Delta = \text{appr}(x, K)$
- write  $\varphi(c)$  increases for  $c \nearrow x$  if there is  $\alpha \in \text{supp } \Delta$  s.th. for all  $c' \in \Delta_\alpha \setminus \{x\}$ :  
 $\varphi(c) > \varphi(c')$  for  $c \nearrow_x c'$   
not necessary if  $\Delta$  is non-trivial  
 where  $\Delta = \text{appr}(x, K)$ .

### Definition

$\Delta$  approximation type over  $(K, v)$

- $f \in K[X]$ :  
 $\Delta$  fixes the value of  $f \iff \exists \alpha \in vK: v(f(c)) = \alpha$  for  $c \nearrow \Delta$   
 Then: this  $\alpha$  is the fixed value  $v(f(c))$  for  $c \nearrow \Delta$
  - $\Delta$  is a transcendental approximation type  
 $\iff \Delta$  fixes the value of every  $f \in K[X]$   
 Otherwise:  $\Delta$  is an algebraic approximation type
    - an associated minimal polynomial for  $\Delta$  is a monic polynomial of minimal degree whose value is not fixed by  $\Delta$
    - the degree of  $\Delta$  is its associated minimal polynomial
- transcendental approximation type: degree  $d = \infty$

### Remark

- An associated minimal polynomial for  $\Delta$  is irreducible
- If  $\text{supp } \Delta \cong vK$ , then the associated minimal polynomial is not unique

### Lemma [KV, Lemma 5.2; Kaplansky, Lemma 4]

$\Gamma$  ordered abelian group,  $\alpha_1, \dots, \alpha_m \in \Gamma$ ,  $\gamma \in \Gamma$  without max element.

$t_1, \dots, t_m \in \mathbb{Z}$ , distinct

Then: ex.  $\beta \in \gamma$ , permutation  $\sigma: \{1, \dots, m\} \rightarrow \{1, \dots, m\}$  s.th. f.a.  $\gamma \in \gamma, \gamma \geq \beta$

$$\alpha_{\sigma(1)} + t_{\sigma(1)} \gamma > \alpha_{\sigma(2)} + t_{\sigma(2)} \gamma > \dots > \alpha_{\sigma(m)} + t_{\sigma(m)} \gamma$$

### Lemma [KV, Lemma 5.2]

$\Delta = \text{appr}(x, K)$  imm. appr. type of degree  $d$  over  $(K, v)$ .

$f \in K[X]$ ,  $\deg(f) \leq d$ ,  $f = \sum c_i X^i$

$f_i := \sum_{j=i}^n \binom{j}{i} c_j X^{j-i}$  formal  $i$ -th derivative of  $f$   $\frac{f^{(i)}}{i!}$

$\beta_i :=$  fixed value  $v(f_i(c))$  for  $c \nearrow x$  ( $\deg(f_i) < \deg(f) \leq d$ )

Then there is  $h \leq \deg f$  such that

$$(5.4) \quad \beta_h + h \cdot v(x-c) < \beta_i + i \cdot v(x-c) \quad \text{for } i+h, i=1, \dots, \deg(f) \\ \text{for } c \nearrow x$$

Hence,

$$(5.5) \quad v(f(x) - f(c)) = \beta_h + h \cdot v(x-c) \quad \text{for } c \neq x$$

Consequently,

$$\begin{cases} v(f(x) - f(c)) > v(f(x)) = v(f(c)) & \text{for } c \neq x, \text{ if } \Delta \text{ fixes the value of } f \\ v(f(x)) > v(f(c)) = \beta_h + h \cdot v(x-c) & \text{for } c \neq x, \text{ if } \Delta \text{ doesn't fix the value of } f \end{cases}$$

**Proof:** Set  $n = \deg(f)$ .

Taylor expansion

$$(5.6) \quad f(x) - f(c) = f_1(c)(x-c) + \dots + f_n(c)(x-c)^n$$

$$v(f_i(c)(x-c)^i) = \frac{\beta_i + i \cdot v(x-c)}{\alpha_i \quad t_i \quad \gamma}$$

$\Delta$  imm. appr. type  $\xrightarrow[\text{v.i.}]{\text{lemma}}$   $\text{supp } \Delta$  has no maximal element

Lemma 5.1. with  $\mathcal{T} = \text{supp } \Delta \subseteq vK$ ,  $\alpha_i = \beta_i$ ,  $t_i = i$ :

There is  $h < \deg(f)$ , s.th. (5.4):

$$\beta_h + h \cdot v(x-c) < \beta_i + i \cdot v(x-c) \quad \text{for } c \neq x, i \neq h$$

This implies (5.5) using (5.6) and ultrametric  $\Delta$ -ineq.

• if  $\Delta$  fixes the value of  $f$ , then assume for contradiction

$$\begin{aligned} v(f(x)) + v(f(c)) &\rightarrow v(f(x) - f(c)) = \min\{v(f(x)), v(f(c))\} \\ \text{for } c \neq x &\quad \text{LHS of (5.5)} \\ &\quad \leftarrow \text{has fixed value for } c \neq x \\ &\quad \downarrow \\ &\quad \text{but } \beta_h + h \cdot v(x-c) \text{ increases for } c \neq x \\ &\quad \text{RHS of (5.5)} \end{aligned}$$

$\Rightarrow \checkmark$

• if  $\Delta$  doesn't fix the value of  $f$ , then  $v(f(c)) \neq v(f(x))$  for  $x \neq c$

$$\Rightarrow v(f(x) - f(c)) = \min\{v(f(x)), v(f(c))\}$$

$$(5.5) = \beta_h + h \cdot v(x-c) \quad \text{constant} \quad \leftarrow \text{the minimum} \Rightarrow \checkmark$$

increases for  $x \neq c$

Lemma 5.2. only for polynomials of degree  $\leq d$ . Now:

Lemma 5.4

Take an immediate algebraic approximation type  $\Delta$ -appr.  $(x, K)$  over  $(K, v)$

and an associated minimal polynomial  $f \in K[X]$  for  $\Delta$ .

Further, take an arbitrary polynomial  $g \in K[X]$  and write

$$g = c_k f^k + \dots + c_1 f + c_0$$

where  $c_i \in K[X]$  with  $\deg(c_i) < \deg(f)$ .

Then there is  $1 \leq m \leq k$ ,  $\beta \in vK$  such that with  $h$  from 5.2

$$v(g(c) - c_0(c)) = v(c_m(c)) + m \cdot v(f(c)) = \beta + m \cdot h \cdot v(x-c) \quad \text{for } c \neq x$$

Consequently,

$$\begin{cases} v(g(x)) = v(g(c)) = v(c_0(c)) = v(c_0(x)) < v(g(c) - c_0(c)) & \text{for } c \neq x \\ & \text{if } \Delta \text{ fixes the value of } g \\ v(g(x)) > v(g(c)) = \beta + m \cdot h \cdot v(x-c) & \text{for } c \neq x \\ & \text{if } \Delta \text{ doesn't fix the value of } g \end{cases}$$

**Proof:** Not in this talk.  $\square$

## § Relative Approximation Degree

$\Delta$  imm. appr. type over  $(K, v)$ ,  $\Delta$ -appr.  $(x, K)$

### Definition

Let  $f \in K[X]$ ,  $\deg(f) \leq \deg \Delta$ . The integer  $h$  from Lemma 5.2. is called the **relative approximation degree of  $f$  in  $x$  over  $K$**  and is denoted by  $h_K(x; f)$ .

### Remark

By Lemma 5.2,

$$1 \leq h_K(x; f) \leq \deg(f)$$

$$\text{and } v(f(x) - f(c)) = \beta_{h_K(x; f)} + h_K(x; f) \cdot v(x - c) \quad \text{for } c \notin x$$

### Observation

$g \in K[X]$ , arbitrary degree.

There are unique  $\beta \in vK$ ,  $k \in \mathbb{Z}_{\geq 0}$  s.th.

$$\text{(\#)} \quad v(g(x) - g(c)) = \beta + k \cdot v(x - c)$$

(This is because  $\Delta$  is immediate and so  $v(x - c)$  takes infinitely many values for  $c \notin x$ )

### Definition

$g \in K[X]$ , arbitrary degree.

The integer  $k$  from (#) is called **relative approximation degree of  $g(x)$  in  $x$** , denoted by  $h_K(x; g)$ .

The value  $\beta$  from (#) is called **relative approximation constant of  $g(x)$  in  $x$** , denoted by  $\beta_K(x; g)$ .

## § Outlook - the proof of 11.1

Situation:

$$(10.1) \quad \begin{cases} (K, v) \text{ a valued field of rank 1} \\ (K(x) | K, v) \text{ an immediate extension s.th. } x \notin K^c \\ \text{appr}(x, K) \text{ is transcendental} \\ y \in K(x)^h \text{ is transcendental over } K \end{cases}$$

Q: What can we say about  $[K(x)^h : K(y)^h]$ ?

(Want:  $[K(x)^h : K(y)^h] = 1$ )

Idea: Define relative approximation degree of  $x$  over  $y$ :

Find polynomial  $f \in K[X]$  s.th.

$$v(y - f(x)) \geq \text{dist}(y, K)$$

↳ smallest initial segment of  $\text{div}(vK)$  containing  $\text{supp appr}(y, K)$

Lemma 10.2: In situation (10.1), such an  $f$  exists

### Definition

$h_K(x; y) := h_K(x; f)$ , the relative appr. degree of  $y$  in  $x$

$\beta_K(x; y) := \beta_K(x; f)$ .

where  $f$  is as in Lemma 10.2.

Lemma 10.3:  $h_K(x:y)$  and  $\beta_K(x:y)$  are well-defined  
i.e. do not depend on  $f$

Theorem 10.7: In situation (10.1),  
 $[K(x)^h : K(y)^h] \leq h_K(x:y)$ . Needs many ingredients

Kuhlmann-Vlahu, Theorem 11.1

$(K,v)$  valued field, rank 1,

$(FK,v)$  immediate function field,  $\text{trdeg}(FK) = 1$

Suppose there is  $x \in F^h \setminus K^c$  with transcendental approximation type over  $K$  such that  $F^h = K(x)^h$

Then there is already some  $y \in F$  such that  $F^h = K(y)^h$   
such that  $F^h = K(y)^h$

Proof idea of 11.1: Find  $y \in F$  such that  $h_K(x:y) = 1$

Then  $[K(x)^h : K(y)^h] \leq h_K(x:y) = 1$  by Thm 10.7.

So  $K(x)^h = K(y)^h$ .

Need dist & stuff from [KV, Chapter 10]