

String diagrams

from control to concurrency and beyond

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Compositionality

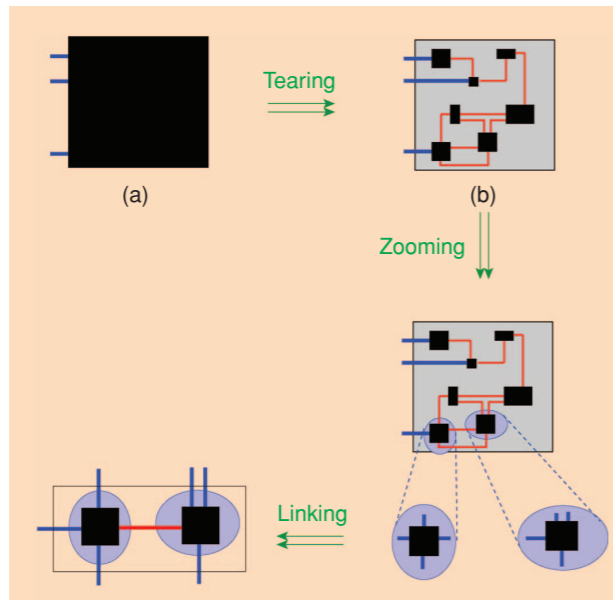


- for “nice” homomorphic translation
 - syntactic operations correspond to natural operations on the semantic domain
 - syntax expressive enough to capture enough of the semantic domain
 - natural notions of semantic equivalence find an axiomatisation in the syntax

Our approach

- in computer science, the tradition is to start with some syntax and study formal semantics as a separate subject
- we think that it is useful to **reverse** the process
 - start with the the algebra of the semantic domain (in CS, control, engineering, science, mathematics, ...)
 - engineer an appropriate syntax to support that algebra

Behavioural control theory



“Thinking of a dynamical system as a behavior, and of **inter-connection as variable sharing**, gets the physics right.”

- Willems’ thesis: abandon causality and functionality (paraphrasing mine)
 - causal thinking is a disease of the brain (Russell, 1912)
 - laws of physics are seldom functional
 - functional modelling is seldom compositional
 - Willems’ tearing procedure produces relational, not functional, behaviours

Compositionality



- What kind of algebra?
 - first order logic, regular logic, relational algebra, datalog, allegories, ...
- What kind of relations?
 - vanilla, additive, linear, affine, ...

Rel_x

- For Willems' intuitions, an appropriate universe seems to be the categorical algebra of the symmetric monoidal category **Rel_x**
 - objects: sets X, Y, Z, \dots
 - arrows: (typed) relations, $R: X \rightarrow Y, S: Y \rightarrow Z$
 - composition: relational composition

$$R ; S = \{ (x,z) \mid \exists y. xRy \wedge ySz \}$$

- monoidal product: $R \times R': X \times X' \rightarrow Y \times Y'$

$$R \times R' = \{ ((x,x'),(y,y')) \mid xRy \wedge x'R'y' \}$$

String diagrams

- diagrammatic syntax for symmetric monoidal categories
- diagrammatic reasoning: the laws of symmetric monoidal categories are baked in to the diagrams

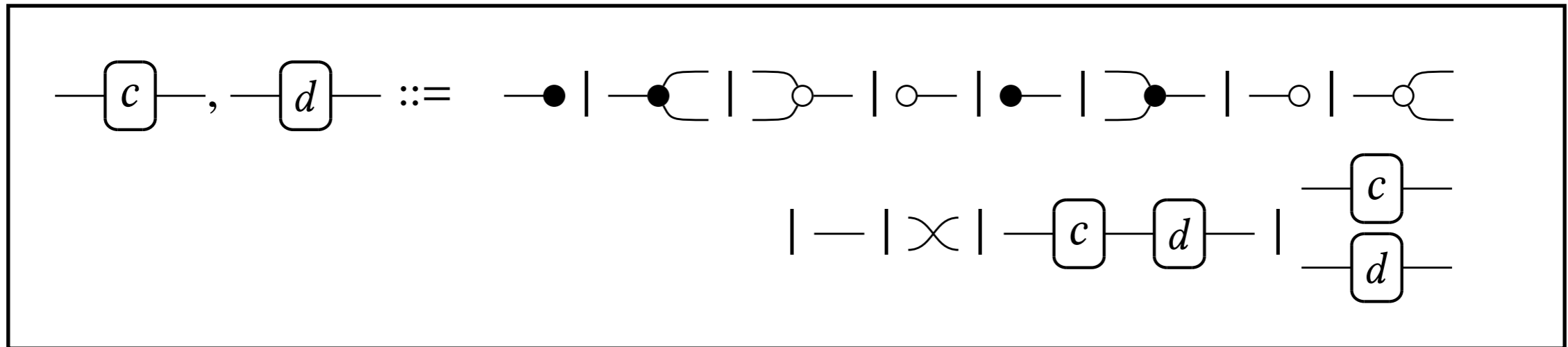
Compositionality



- syntax expressive enough?
- axiomatisations?

Graphical Linear Algebra

- String diagrams generated by the following syntax



String diagrams

→
monoidal functor

LinRel₀

Sound and fully complete axiomatisation - the theory of **IH** (Interacting Hopf algebras)

(Bonchi, S., Zanasi, Interacting Hopf Algebras, 2014)

Signal flow graphs

- The **IH** construction is parametric wrt any PID
- Starting with $\mathbf{R}[x]$ we get linear relations over its field of fractions $\mathbf{R}(x)$
- This yields a sound and complete equational system for reasoning about signal flow graphs: models of computation that compute solutions of rational functions

F. Bonchi, P. Sobociński and F. Zanasi, "[Full Abstraction for Signal Flow Graphs](#)", *In Principles of Programming Languages, POPL`15*

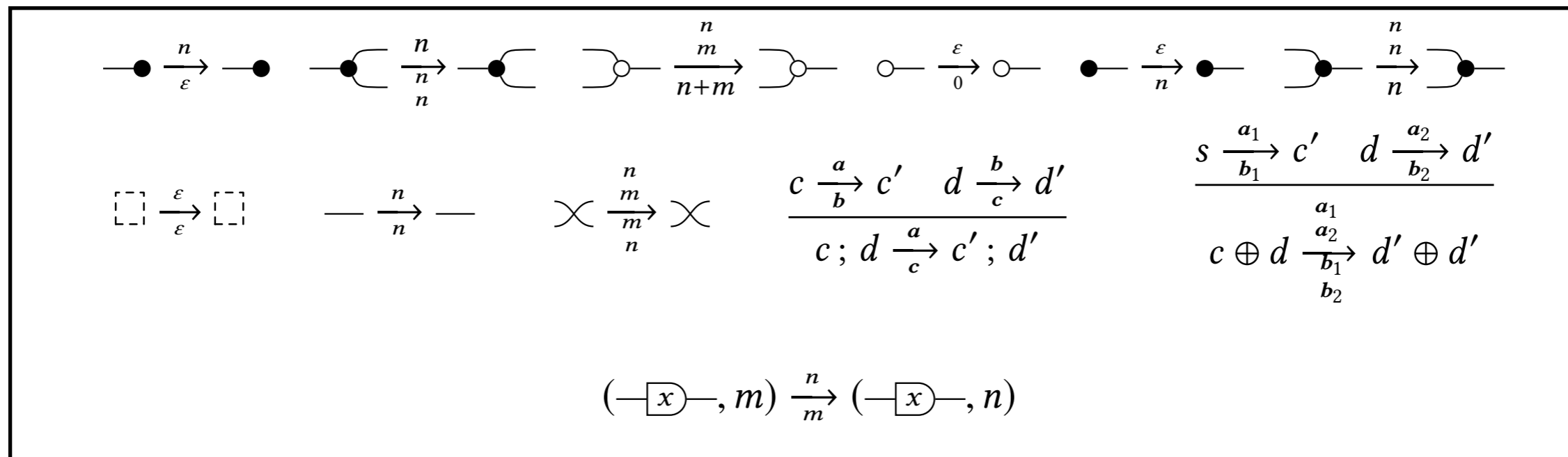
F. Bonchi, P. Sobociński and F. Zanasi, "[The Calculus of Signal Flow Diagrams I: Linear Relations on Streams](#)", *Inf Comput*

B. Fong, P. Rapisarda and P. Sobociński, "[A categorical approach to open and interconnected dynamical systems](#)", *LICS`16*

F. Bonchi, J. Holland, D. Pavlovic and P. Sobociński, "[Refinement for signal flow graphs](#)", *CONCUR`17*

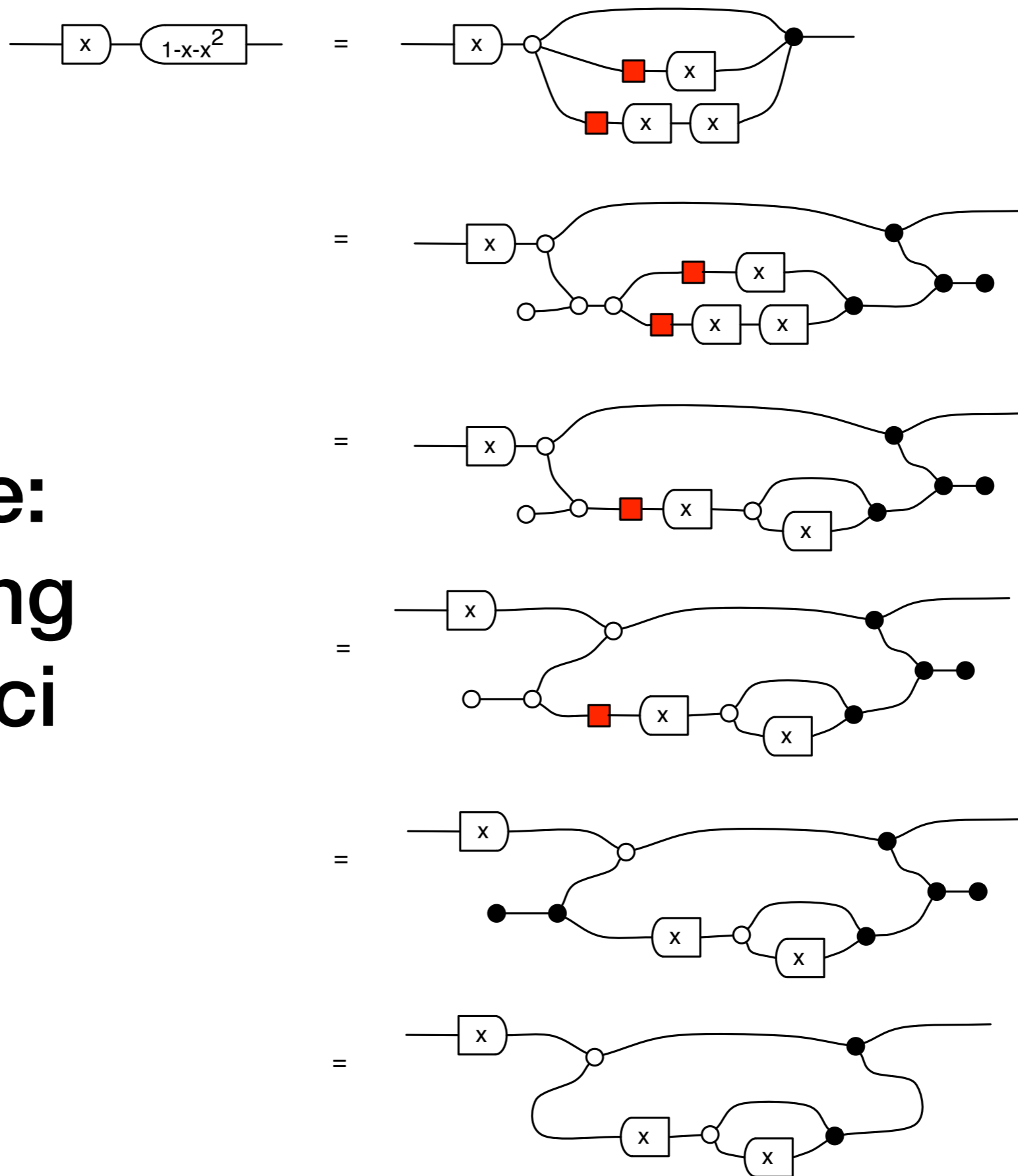
The operational view

- The work on signal flow graphs emphasises the importance of the **operational view**



- For signal flow graphs, the signals come from a field, typically **R** or **Q**

Example: computing Fibonacci



Graphical Diophantine Algebra

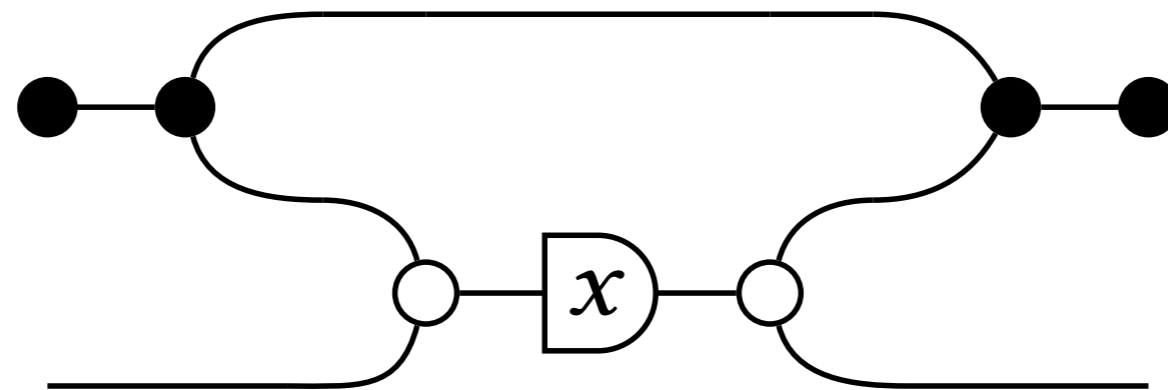
- **Definition.** An additive relation of type $k \rightarrow l$ is a subset $R \subseteq \mathbf{N}^k \times \mathbf{N}^l$ s.t. $(0,0) \in \mathbf{R}$ and, if $(a,b), (a',b') \in \mathbf{R}$ then $(a+a',b+b') \in \mathbf{R}$
- An additive relation is f.g. if we can find a finite basis: i.e. every element can be expressed as a sum of basis elements
- These form a prop **AddRel** as a subprop of **Rel_x**
 - proving f.g. additive relations are closed under composition is a cute application of Dickson's Lemma



Same syntax as before, and.... sound and fully complete axiomatisation

From control to concurrency

- For linear relations, adding state yielded a compositional account of signal flow graphs
- For additive relations, adding state yields a compositional account of Petri nets



Graphical Affine Algebra

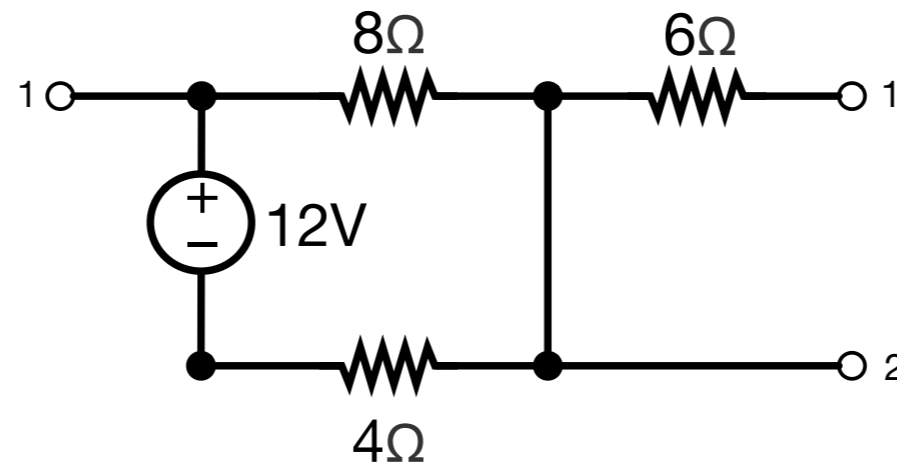


- The usual syntax extended with \vdash that “outputs 1”

Two sound and complete axiomatisations.

Fun application: electrical circuits

- Let's go back to the \mathbf{R} world. We will use Graphical Affine Algebra as a sound and complete diagrammatic proof system for open circuits like:



Compiling to string diagrams

$$\mathcal{I} \left(\text{---} \begin{array}{|c} \text{---} \\ \text{---} \end{array} \right) = \text{---} \begin{array}{|c} \bullet \\ \circ \end{array} \text{---}$$

$$\mathcal{I} \left(\begin{array}{|c} \text{---} \\ \text{---} \end{array} \text{---} \right) = \text{---} \begin{array}{|c} \circ \\ \bullet \end{array} \text{---}$$

$$\mathcal{I} (\bullet \text{---}) = \begin{array}{|c} \bullet \\ \circ \end{array} \text{---}$$

$$\mathcal{I} (\text{---} \bullet) = \text{---} \begin{array}{|c} \bullet \\ \circ \end{array}$$

$$\mathcal{I} \left(\begin{array}{c} k \\ \text{---} \text{---} \end{array} \right) = \text{---} \begin{array}{|c} \circ \\ \bullet \end{array} \text{---} \text{---} \begin{array}{|c} \bullet \\ \circ \end{array} \text{---}$$

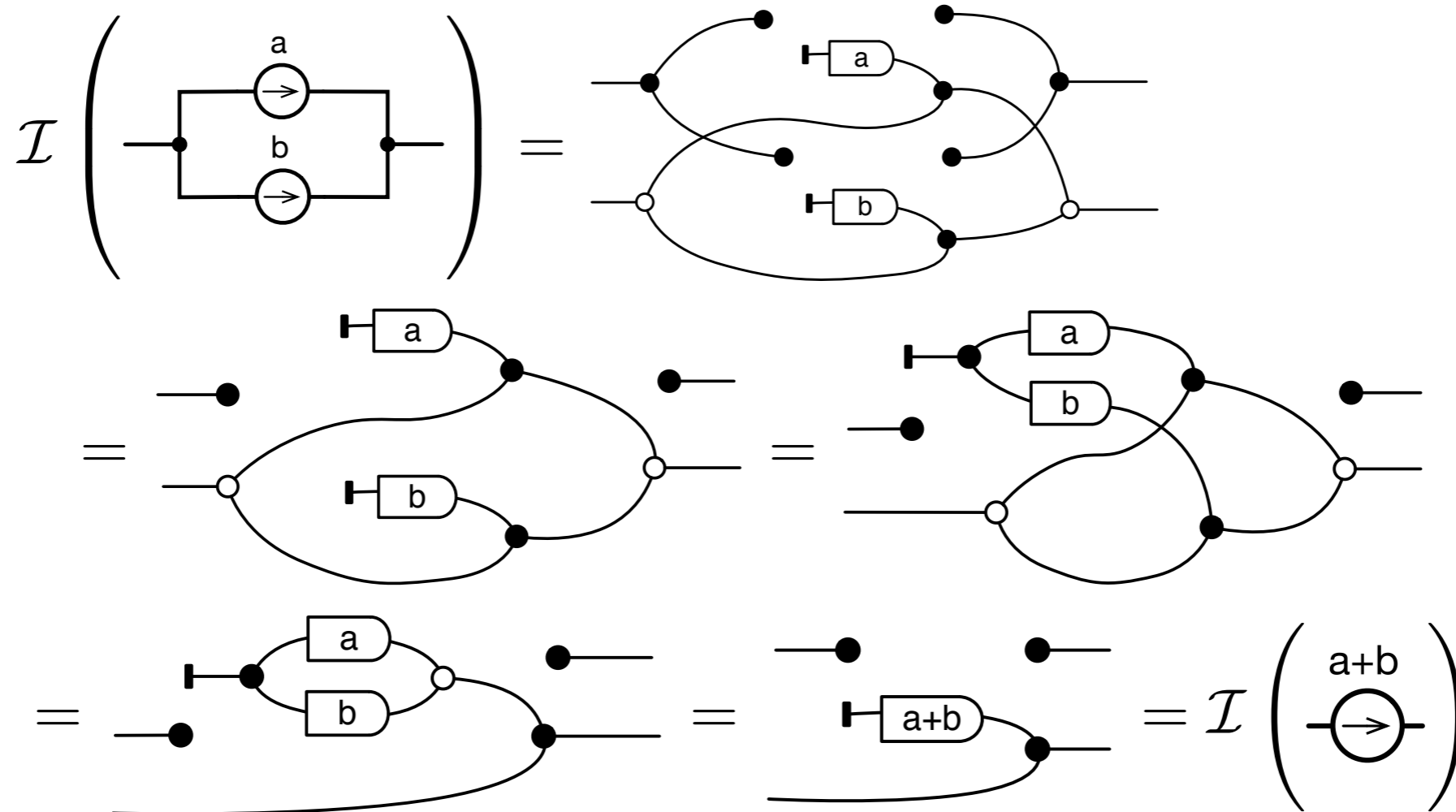
$$\mathcal{I} \left(\begin{array}{c} k \\ \text{---} \oplus \text{---} \end{array} \right) = \text{---} \begin{array}{|c} \circ \\ \bullet \end{array} \text{---} \text{---} \begin{array}{|c} \bullet \\ \circ \end{array} \text{---}$$

$$\mathcal{I} \left(\begin{array}{c} k \\ \text{---} \otimes \text{---} \end{array} \right) = \text{---} \begin{array}{|c} \bullet \\ \circ \end{array} \text{---} \text{---} \begin{array}{|c} \bullet \\ \circ \end{array} \text{---}$$

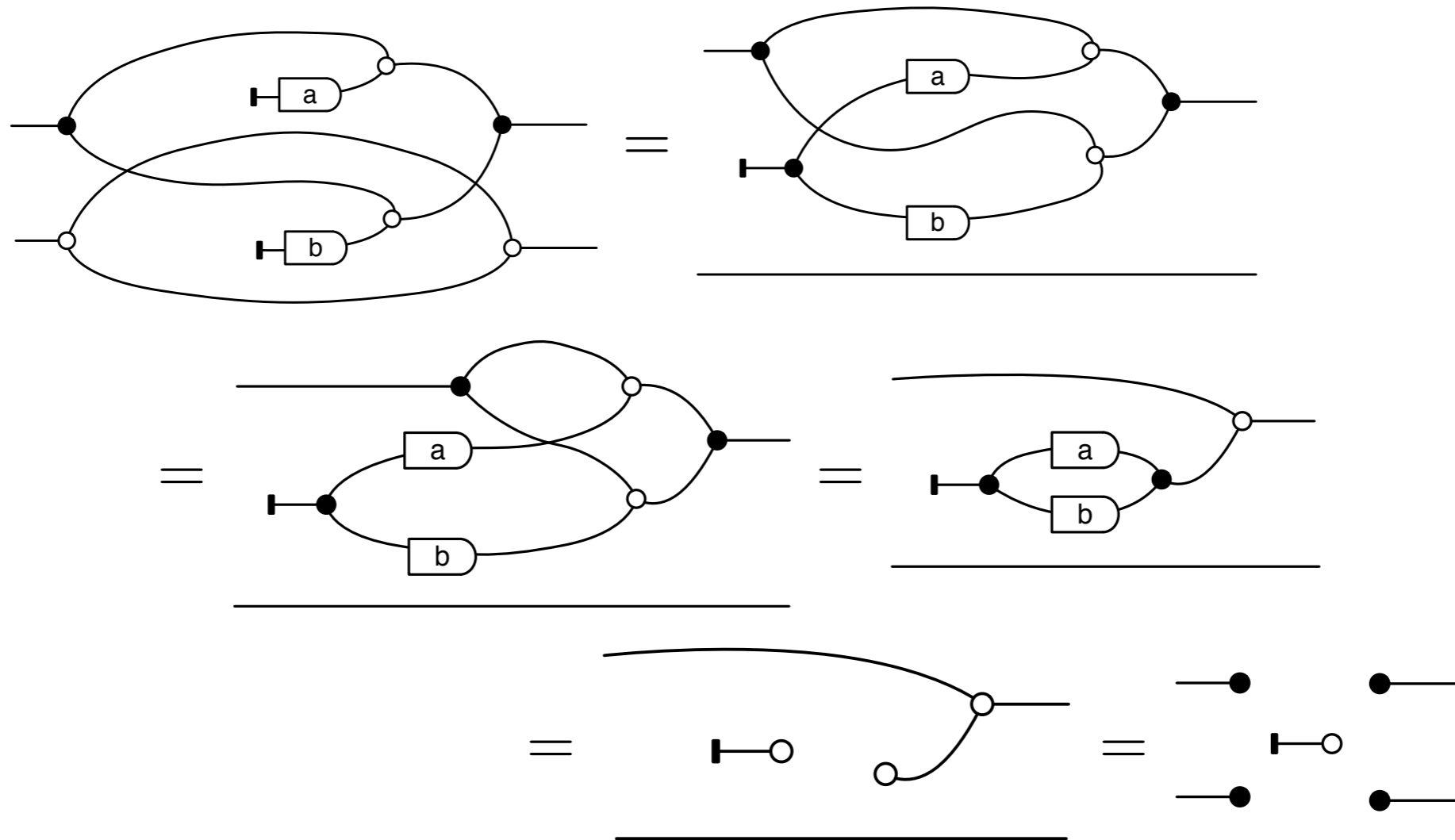
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Current sources in parallel are additive



Voltage sources in parallel are “illegal”



- inspired by process algebra - operational playing around leads to equations leads to denotations
- unlike process algebra, we are not reinventing the algebraic wheel: the basic operations for composing process are those of monoidal categories
- what most surprises me is **robustness**.
 - on the semantic side, the mathematics changes drastically
 - equationally, in terms of the string diagrams, we change some basic interaction of GLA primitives

Compositional systems and methods



- new compositionality group at Taltech: applications of category theory to concurrency, control, game theory, engineering, machine learning, ...
- come and visit!!
- SYCO 7 in Tallinn - March 30-31, 2020 - save the date!

