

# **Component-based semantics**

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# **PlanCompS: Programming Language Components and Specifications**

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# Main idea

## A component-based framework

- ▶ to support design, specification, implementation of programming and domain-specific languages

## with some novel aspects:

- ▶ an open-ended collection of highly reusable language components: **funcons** (fundamental constructs)
- ▶ each funcon is to be specified **independently**
- ▶ a **digital library** of language and funcon specifications

# Background

## Programming languages

- ▶ **many hundreds of them**

- **older:** Fortran, Cobol, Ada, C++, Scheme, Haskell, ...
- **newer:** Java, C#, Python, Ruby, OCaml, VBScript, ...
- **domain-specific:** PHP, Javascript, ...

- ▶ **continually evolving**

- **new versions** of existing languages
- **new languages**

# Background

## Programming language specifications

- ▶ **program syntax** – formal

- lexical symbols
- phrase structure

- ▶ **program semantics** – informal

- static (compile-time) behaviour
- dynamic (run-time) behaviour

# Background

Current **pragmatic problems** with formal semantics concerning:

- ▶ **reuse**
- ▶ **co-evolution**
- ▶ **tool support**

when specifying *major* or *rapidly-evolving* languages

# PLANCOMPS

Programming  
Languages  
(incl. DSLs)

C#

Java ...

translation

Components  
and their  
Specifications

reusable

● ● ● ● ● ● ● ● ● ● ● ● ● ● ● ...

fundamental constructs  
'funcons'

# Component-based specifications

**Syntax** (concrete, abstract)

- ▶ BNF + regular expressions

**Semantics** (static, dynamic)

- ▶ context-free translation to funcons

**Funcons**

- ▶ modular SOS rules
- ▶ modular bisimilarity theory



# OCaml Light

## **OCaml** (Objective Caml, F#)

- ▶ a popular OO functional language

## **Caml Light**

- ▶ a pedagogical sublanguage of Caml

## **OCaml Light**

- ▶ defined in SOS using Ott by Scott Owens

# **OCaml Light syntax**

# OCaml Light syntax

## Outlines

=====

### Lexical non-terminals

-----

...

### Names

-----

I : ident ::= lowercase\_ident  
| capitalized\_ident

VN : value\_name ::= lowercase\_ident  
| '(' operator\_name ')'

ON : operator\_name ::= prefix\_symbol  
| infix\_op

infix\_op ::= infix\_symbol  
| '\*' | '=' | 'or' | '&  
| ':=' | 'mod' | 'land' | 'lor'  
| 'lxor' | 'asr' | 'lsl' | 'lsr'

IO : infix\_op \ ( '&&' | '&' | '||' | 'or' )

K : constr ::= capitalized\_ident

TK : typeconstr ::= lowercase\_ident

F : field ::= lowercase\_ident

### Type expressions

-----

T : typexpr ::= ...

### Constants

-----

C : constant ::= integer\_literal  
| float\_literal  
| char\_literal  
| string\_literal  
| constr  
| 'false'  
| 'true'  
| '[' ']'  
| '(' ')'

## Expressions

-----

E : expr ::= value\_name  
| constant  
| '(' expr ')'  
| 'begin' expr 'end'  
| '(' expr ':' typexpr ')'  
| expr ( ',' expr )+  
| constr expr  
| expr '::' expr  
| '[' expr ( ';' expr )\* ']'  
| '{' field '=' expr  
  ( ';' field '=' expr )\* '}'  
| '{' expr 'with'  
  field '=' expr  
  ( ';' field '=' expr )\* '}'  
| expr expr+  
| prefix\_symbol expr  
| ( '+' | '-' | '+.' | '-.' ) expr  
| expr infix\_op expr  
| expr '.' field  
| 'if' expr 'then' expr ( 'else' expr )?  
| 'while' expr 'do' expr 'done'  
| 'for' value\_name '=' expr  
  ( 'to' | 'downto' ) expr  
  'do' expr 'done'  
| expr ';' expr  
| 'match' expr 'with' pattern\_matching  
| 'function' pattern\_matching  
| 'fun' pattern+ '->' expr  
| 'try' expr 'with' pattern\_matching  
| 'let' 'rec'? let\_binding  
  ( 'and' let\_binding )\*  
  'in' expr  
| 'assert' expr

## Patterns

-----

P : pattern ::= value\_name  
| '\_'  
| constant  
| pattern 'as' value\_name  
| '(' pattern ')'  
| '(' pattern ':' typexpr ')'  
| pattern '|' pattern  
| constr pattern  
| pattern ( ',' pattern )+  
| '{' field '=' pattern  
  ( ';' field '=' pattern )\* '}'  
| '[' pattern ( ';' pattern )\* ']'  
| pattern '::' pattern

## Pattern matching

-----

PM : pattern\_matching ::=  
  '|'? pattern '->' expr  
  ( '|' pattern '->' expr )\*

## Declarations and definitions

-----

LB : let\_binding ::=  
  pattern '=' expr  
  | value\_name pattern+ '=' expr  
  | value\_name ':' typexpr '=' expr  
  | value\_name pattern+ ':' typexpr '=' expr

TD : type\_definition ::= ...

ED : exception\_definition ::= ...

D : definition ::=  
  'let' 'rec'? let\_binding  
  ( 'and' let\_binding )\*  
  | type\_definition  
  | exception\_definition

TP : toplevel\_phrase ::= definition | expr

TI : toplevel\_input ::= toplevel\_phrase\* ';;'

## Programs

-----

PR: program ::= toplevel\_input+

# OCaml Light syntax

## Lexical non-terminals

---

lowercase\_ident  
capitalized\_ident  
prefix\_symbol  
infix\_symbol  
integer\_literal  
float\_literal  
char\_literal  
string\_literal

## Names

---

ident  
operator\_name  
infix\_op  
constr  
typeconstr  
field

## Type expressions

---

typexpr

## Constants

---

constant

## Expressions

---

expr

## Patterns

---

pattern

## Pattern matching

---

pattern\_matching

## Declarations and definitions

---

let\_binding  
type\_definition  
typedef  
type\_params  
type\_param  
constr\_decl  
field\_decl  
exception\_definition  
definition  
toplevel\_phrase  
toplevel\_input  
program

# OCaml Light syntax

Lexical non-terminals

.....

Names

-----

```
I : ident          ::= lowercase_ident | capitalized_ident
VN : value_name    ::= lowercase_ident | '(' operator_name ') '
ON : operator_name ::= prefix_symbol   | infix_op
```

```
infix_op          ::= infix_symbol
                  | '*'      | '='      | 'or'      | '&'
                  | ':='     | 'mod'   | 'land'   | 'lor'
                  | '\xor'  | 'asr'  | '\sl'   | '\sr'
```

```
I0 : infix_op \ ( '&&' | '&' | '||' | 'or' )
```

```
K : constr        ::= capitalized_ident
F : field         ::= lowercase_ident
```

# OCaml Light syntax

Type expressions

-----

T : typexpr ::= ...

Constants

-----

C : constant ::= integer\_literal  
                  | float\_literal  
                  | char\_literal  
                  | string\_literal  
                  | constr  
                  | 'false'  
                  | 'true'  
                  | '[' ' ]'  
                  | '(' ' )'

# OCaml Light syntax

## Expressions

```
-----  
E : expr ::= value_name  
      | constant  
      | '(' expr ')'  
      | 'begin' expr 'end'  
      | '(' expr ':' typexpr ')'  
      | expr ( ',' expr )+  
      | constr expr  
      | expr '::' expr  
      | '[' expr ( ';' expr )* ']'  
      | '{' field '=' expr ( ';' field '=' expr )* '}'  
      | '{' expr 'with'  
          field '=' expr ( ';' field '=' expr )* '}'  
      | expr expr+  
      | prefix_symbol expr  
      | ( '+' | '-' | '+.' | '-.' ) expr  
      | expr infix_op expr  
      | ...
```

# OCaml Light syntax

## Expressions

```
-----  
E : expr ::= ...  
| expr '.' field  
| 'if' expr 'then' expr ( 'else' expr )?  
| 'while' expr 'do' expr 'done'  
| 'for' value_name '=' expr  
|   ( 'to' | 'downto' ) expr  
|   'do' expr 'done'  
| expr ';' expr  
| 'match' expr 'with' pattern_matching  
| 'function' pattern_matching  
| 'fun' pattern+ '->' expr  
| 'try' expr 'with' pattern_matching  
| 'let' 'rec'? let_binding  
|   ( 'and' let_binding )*  
|   'in' expr  
| 'assert' expr
```



# OCaml Light syntax

## Patterns

```
-----  
P : pattern ::= value_name  
      | '_'  
      | constant  
      | pattern 'as' value_name  
      | '(' pattern ')'  
      | '(' pattern ':' typexpr ')'  
      | pattern '|' pattern  
      | constr pattern  
      | pattern ( ',' pattern )+  
      | '{' field '=' pattern ( ';' field '=' pattern )* '}'  
      | '[' pattern ( ';' pattern )* ']'  
      | pattern '::' pattern
```

## Pattern matching

```
-----  
PM : pattern_matching ::=  
      '|'? pattern '->' expr ( '|' pattern '->' expr )*
```

# OCaml Light syntax

Declarations and definitions

---

```
LB : let_binding ::= pattern '=' expr
    | value_name pattern+ '=' expr
    | value_name ':' typexpr '=' expr
    | value_name pattern+ ':' typexpr '=' expr

TD : type_definition ::= ...

ED : exception_definition ::= ...

D : definition ::= 'let' 'rec'? let_binding ( 'and' let_binding )*
    | type_definition
    | exception_definition

TP : toplevel_phrase ::= definition | expr

TI : toplevel_input ::= toplevel_phrase* ';;'
```

Programs

---

```
PR: program ::= toplevel_input+
```

# **OCaml Light Semantics**

# Translation to funcons

<i>id</i> [[ <i>VN</i> ]]	: <i>id</i>	<i>VN</i> : value_name
<i>value</i> [[ <i>C</i> ]]	: <i>value</i>	<i>C</i> : constant
<i>expr</i> [[ <i>E</i> ]]	: <i>expr</i>	<i>E</i> : expr
<i>patt</i> [[ <i>P</i> ]]	: <i>patt</i>	<i>P</i> : pattern
<i>abs</i> [[ <i>PM</i> ]]	: <i>abs</i>	<i>PM</i> : pattern_matching
<i>decl</i> [[ <i>LB</i> ]]	: <i>decl</i>	<i>LB</i> : let_binding
<i>decl</i> [[ <i>D</i> ]]	: <i>decl</i>	<i>D</i> : definition
<i>decl</i> [[ <i>TP</i> ]]	: <i>decl</i>	<i>TP</i> : toplevel_phrase
<i>decl</i> [[ <i>TI</i> ]]	: <i>decl</i>	<i>TI</i> : toplevel_input
<i>decl</i> [[ <i>PR</i> ]]	: <i>decl</i>	<i>PR</i> : program

# Funcons

# Computation types

Type **computes** ( $X$ )

Type **comm** = **computes** (**skip**)

Type **decl** = **computes** (**env**)

Type **expr** = **computes** (**value**)

Type **depends** ( $X, Y$ )

Type **abs** = **depends** (**passable**, **value**)

Type **patt** = **depends** (**value**, **env**)

# Value types

**atom** < value  
**boolean** < value  
**int** < value  
**float** < value  
**char** < value  
**string** < value  
**tuple** < value  
**record** < value  
**variant** < value  
**function** < value  
**variable** < value  
**skip** < value

**passable** = value  
**throwable** = value  
**bindable** = value  
**storable** = value





# *Preliminary* tool support

IDE :

- ▶ **Spoofax**

Parsing, translation to funcons :

- ▶ **ASF+SDF**

Funcon interpretation :

- ▶ **Prolog**





# **OCaml Light Semantics**

# Translation to funcons

<i>id</i> [[ <i>VN</i> ]]	: <b>id</b>	<i>VN</i> :	value_name
<i>value</i> [[ <i>C</i> ]]	: <b>value</b>	<i>C</i> :	constant
<i>expr</i> [[ <i>E</i> ]]	: <b>expr</b>	<i>E</i> :	expr
<i>patt</i> [[ <i>P</i> ]]	: <b>patt</b>	<i>P</i> :	pattern
<i>abs</i> [[ <i>PM</i> ]]	: <b>abs</b>	<i>PM</i> :	pattern_matching
<i>decl</i> [[ <i>LB</i> ]]	: <b>decl</b>	<i>LB</i> :	let_binding
<i>decl</i> [[ <i>D</i> ]]	: <b>decl</b>	<i>D</i> :	definition
<i>decl</i> [[ <i>TP</i> ]]	: <b>decl</b>	<i>TP</i> :	toplevel_phrase
<i>decl</i> [[ <i>TI</i> ]]	: <b>decl</b>	<i>TI</i> :	toplevel_input
<i>decl</i> [[ <i>PR</i> ]]	: <b>decl</b>	<i>PR</i> :	program

# Value and operator names

**SYNTAX**       $VN : \text{value\_name}$

**SEMANTICS**     $id[[ VN ]] : id$

\*  $\text{value\_name} ::= \text{lowercase\_ident}$

$id[[ LI ]] = id(LI)$

\*  $\text{value\_name} ::= '(' \text{operator\_name} ')'$

$id[[ ( ON ) ]] = id(ON)$

**SYNTAX**       $ON : \text{operator\_name}$

\*  $\text{operator\_name} ::= \text{prefix\_symbol}$

\*  $\text{operator\_name} ::= \text{infix\_op}$

# Literal constants

**SYNTAX**       $C$  : constant

**SEMANTICS**     $value$ [[  $C$  ]] :  $value$

\* constant ::= integer\_literal

$value$ [[  $IL$  ]] =  $IL$

int < value

\* constant ::= float\_literal

$value$ [[  $FL$  ]] =  $FL$

float < value

...

char < value

...

string < value

# Constructors and booleans

**SYNTAX**       $C$  : constant

**SEMANTICS**     $value$ [[  $C$  ]] :  $value$

\* constant ::= constr

$value$ [[  $K$  ]] =  $K$

atom < value

\* constant ::= 'false' | 'true'

$value$ [[ **false** ]] = **false**

$value$ [[ **true** ]] = **true**

boolean < value



# Empty list and tuple

**SYNTAX**       $C$  : constant

**SEMANTICS**     $value$  [[  $C$  ]] :  $value$

\* constant ::= '[' ' ]'

$value$  [[ [ ] ]] = '[]'

atom < value

\* constant ::= '(' ' )'

$value$  [[ ( ) ]] = **tuple\_empty**

tuple < value

# Expressions

# Simple expressions

**SYNTAX**       $E : \text{expr}$

**SEMANTICS**     $\text{expr}[[ E ]] : \text{expr}$

\*  $\text{expr} ::= \text{value\_name}$

$\text{expr}[[ VN ]] = \text{bound\_value}(\text{id}[[ VN ]])$

\*  $\text{expr} ::= \text{constant}$

$\text{expr}[[ C ]] = \text{value}[[ C ]]$

# Grouped expressions

\* `expr ::= '(' expr ')'`

**`expr`**[[ ( *E* ) ]] = **`expr`**[[ *E* ]]

\* `expr ::= 'begin' expr 'end'`

**`expr`**[[ **begin** *E* **end** ]] = **`expr`**[[ *E* ]]

\* `expr ::= '(' expr ':' typexpr ')'`

**`expr`**[[ ( *E* : *T* ) ]] = **`expr`**[[ *E* ]]

*// for dynamic semantics only*

# Tuple expressions

\*  $\text{expr} ::= \text{expr} ( ', ' \text{expr} )^+$

$\text{expr} [[ E1 , E2 \dots ]] =$   
 $\text{tuple\_prefix}(\text{expr} [[ E1 ]],$   
 $\text{expr\_tuple} [[ E2 \dots ]])$

tuple < value

$\dots : ( ', ' \text{expr} )^*$

**SEMANTICS**  $\text{expr\_tuple} [[ E \dots ]] : \text{expr}$

–  $\text{expr\_tuple} [[ E1 ]] =$   
 $\text{tuple\_prefix}(\text{expr} [[ E1 ]], \text{tuple\_empty})$

–  $\text{expr\_tuple} [[ E1 , E2 \dots ]] =$   
 $\text{tuple\_prefix}(\text{expr} [[ E1 ]],$   
 $\text{expr\_tuple} [[ E2 \dots ]])$

# Constructors and lists

\*  $\text{expr} ::= \text{constr expr}$

$\text{expr} [[ K E ]] = \text{variant}(K, \text{expr} [[ E ]])$

variant < value

\*  $\text{expr} ::= \text{expr} '::' \text{expr}$

$\text{expr} [[ E1 :: E2 ]] =$   
 $\text{variant}('::', \text{tuple2}(\text{expr} [[ E1 ]],$   
 $\text{expr} [[ E2 ]]))$

variant < value

\*  $\text{expr} ::= '[' \text{expr} ( ';' \text{expr} )^* '['$

$\text{expr} [[ [ E1 ] ]] = \text{expr} [[ E1 :: [ ] ]]$

$\text{expr} [[ [ E1 ; E2 \dots ] ]] = \text{expr} [[ E1 :: [ E2 \dots ] ]]$

$\dots : ( ';' \text{expr} )^*$

# Record expressions

\*  $\text{expr} ::= \{ \text{field} '=' \text{expr} \}$   
           $( ';' \text{field} '=' \text{expr} )^*$

$\text{expr}[[ \{ F1 = E1 \} ]]$  = **record1**( $F1$ ,  $\text{expr}[[ E1 ]]$ )

$\text{expr}[[ \{ F1 = E1 ; F2 = E2 \dots \} ]]$  =  
      **record\_union**( $\text{expr}[[ \{ F1 = E1 \} ]]$ ,  
                       $\text{expr}[[ \{ F2 = E2 \dots \} ]]$ )

record < value

$\dots : ( ';' \text{field} '=' \text{expr} )^*$

\*  $\text{expr} ::= \text{expr} '.' \text{field}$

$\text{expr}[[ E . F ]]$  = **record\_select**( $\text{expr}[[ E1 ]]$ ,  $F$ )

record < value

# Record expressions

```
* expr ::= '{' expr 'with' field '=' expr  
          ( ';' field '=' expr )* '}'
```

```
expr[[ { E with F1 = E1 } ]] =  
  record_over(expr[[ { F1 = E1 } ]], expr[[ E ]])
```

```
expr[[ { E with F1 = E1 ; F2 = E2 ... } ]] =  
  expr[[ { { E with F1 = E1 } with F2 = E2 ... } ]]
```

record < value

```
... : ( ';' field '=' expr)*
```



# Function application

\*  $\text{expr} ::= \text{expr expr}^+$

$\text{expr} [[ E1 E2 ] ] = \text{apply\_function}(\text{expr} [[ E1 ] ], \text{expr} [[ E2 ] ])$

$\text{expr} [[ E1 E2 E3 \dots ] ] = \text{expr} [[ (E1 E2) E3 \dots ] ]$

function < value

$\dots : \text{expr}^*$

# Prefix operators

\* `expr ::= prefix_op expr`

`expr` [[ *P0* *E1* ]] = `expr` [[ ( *P0* ) *E1* ]]

\* `expr ::= ( '+' | '-' | '+.' | '-.' ) expr`

`expr` [[ + *E1* ]] = `int_plus`(0, `expr` [[ *E1* ]])

`expr` [[ - *E1* ]] = `int_minus`(0, `expr` [[ *E1* ]])

`int` < value

`expr` [[ +. *E1* ]] = `float_plus`(0.0, `expr` [[ *E1* ]])

`expr` [[ -. *E1* ]] = `float_minus`(0.0, `expr` [[ *E1* ]])

`float` < value

# Infix operators

\*  $\text{expr} ::= \text{expr infix\_op expr}$

$\text{expr} [[ E1 \ I0 \ E2 ]] = \text{expr} [[ ( \ I0 \ ) \ E1 \ E2 ]]$

\*  $I0 : \text{infix\_op} \setminus ( \ \&\&\ | \ \&\ | \ \|\|\ | \ \text{or}\ )$

$\text{expr} [[ E1 \ \&\&\ E2 ]] = \text{if\_true}(\text{expr} [[ E1 ]],$   
 $\text{expr} [[ E2 ]],$   
 $\text{false})$

$\text{expr} [[ E1 \ \&\ E2 ]] = \text{expr} [[ E1 \ \&\&\ E2 ]]$

$\text{expr} [[ E1 \ \|\|\ E2 ]] = \text{if\_true}(\text{expr} [[ E1 ]],$   
 $\text{true},$   
 $\text{expr} [[ E2 ]])$

$\text{expr} [[ E1 \ \text{or}\ E2 ]] = \text{expr} [[ E1 \ \|\|\ E2 ]]$

boolean < value

# Conditionals and loops

\* `expr ::= 'if' expr 'then' expr ( 'else' expr )?`

`expr`[[ **if** `E1` **then** `E2` ]] =  
`expr`[[ **if** `E1` **then** `E2` **else** ( ) ]]

`expr`[[ **if** `E1` **then** `E2` **else** `E3` ]] =  
**if\_true**(`expr`[[ `E1` ]], `expr`[[ `E2` ]], `expr`[[ `E3` ]])

boolean < value

\* `expr ::= 'while' expr 'do' expr 'done'`

`expr`[[ **while** `E1` **do** `E2` **done** ]] =  
**seq**(**while\_true**(`expr`[[ `E1` ]], **effect**(`expr`[[ `E2` ]])),  
**tuple\_empty**)

boolean < value

# For-loops

```
* expr ::= 'for' value_name '=' expr  
        ( 'to' | 'downto' ) expr  
        'do' expr 'done'
```

```
expr[[ for VN = E1 to E2 do E3 done ]] =  
    apply_to_each(abs(bind(id[[ VN ]]),  
                  effect(expr[[ E3 ]])),  
    int_closed_interval(expr[[ E1 ]],  
                        expr[[ E2 ]]))
```

```
expr[[ for VN = E1 downto E2 do E3 done ]] =  
    apply_to_each(abs(bind(id[[ VN ]]),  
                  effect(expr[[ E3 ]])),  
    list_reverse(int_closed_interval(expr[[ E2 ]],  
                                      expr[[ E1 ]]))
```

```
int < passable
```

# Sequential expressions

\*  $\text{expr} ::= \text{expr} \text{ ; } \text{expr}$

$\text{expr} [[ E1 \ ; \ E2 \ ]] = \text{seq}(\text{effect}(\text{expr} [[ E1 \ ]]), \text{expr} [[ E2 \ ]])$

done < value

# Match expressions

\* `expr ::= 'match' expr 'with' pattern_matching`

```
expr [[ match E with PM ]] =  
    apply(prefer_over(abs [[ PM ]],  
                    abs(any, throw( 'Match_failure' ))),  
    expr [[ E ]])
```

'Match\_failure' : throwable

# Function abstraction

```
* expr ::= 'function' pattern_matching
```

```
expr[[ function PM ]] =  
    close function(prefer_over(abs[[ PM ]],  
                                abs(any, throw('Match_failure'))))
```

```
'Match_failure' : throwable
```

```
* expr ::= 'fun' pattern+ '->' expr
```

```
expr[[ fun P1 -> E ]] =  
    expr[[ function P1 -> E ]]
```

```
expr[[ fun P1 P2 ... -> E ]] =  
    expr[[ fun P1 -> fun P2 ... -> E ) ]]
```

```
... : pattern*
```



# Exception handling

\* `expr ::= 'try' expr 'with' pattern_matching`

```
expr [[ try E with PM ]] =  
  catch(expr [[ E ]],  
        prefer_over(abs [[ PM ]], abs(throw(given))))
```

# Let-expressions

```
* expr ::= 'let' 'rec'? let_binding  
          ( 'and' let_binding )* 'in' expr
```

```
expr[[ let LB ... in E ]] =  
  scope(decl[[ LB ... ]], expr[[ E ]])
```

```
expr[[ let rec LB ... in E ]] =  
  scope(recursive(decl[[ LB ... ]]), expr[[ E ]])
```

```
... : ('and' let_binding)*
```

# Assertions

\* `expr ::= 'assert' expr`

```
expr [[ assert E ]] =  
    if_true(expr [[ E ]],  
            tuple_empty,  
            throw( 'Assert_failure' ))
```

`boolean < value`

# Patterns

# Simple patterns

**SYNTAX**       $P$  : pattern

**SEMANTICS**     $\mathit{patt}[[ P ]]$  :  $\mathit{patt}$

\* pattern ::= value\_name

$\mathit{patt}[[ VN ]]$  =  $\mathit{bind}(\mathit{id}[[ VN ]])$

\* pattern ::= '\_'

$\mathit{patt}[[ _ ]]$  =  $\mathit{any}$

\* pattern ::= constant

$\mathit{patt}[[ C ]]$  =  $\mathit{only}(\mathit{value}[[ C ]])$

# Composite patterns

\* pattern ::= pattern 'as' value\_name

*patt*[[ *P* as *VN* ]] = **map\_union**(*patt*[[ *P* ]],  
*patt*[[ *VN* ]])

\* pattern ::= pattern '|' pattern

*patt*[[ *P1* | *P2* ]] = **prefer\_over**(*patt*[[ *P1* ]],  
*patt*[[ *P2* ]])

# Constructor patterns

\* pattern ::= constr pattern

```
patt[[ K P ]] =  
  invert variant(only(K), patt[[ P ]])
```

variant < value

\* pattern ::= pattern ':::' pattern

```
patt[[ P1 ::: P2 ]] =  
  invert variant(only(':::'),  
    invert tuple2(patt[[ P1 ]],  
      patt[[ P2 ]]))
```

# List patterns

\* pattern ::= '[' pattern ( ';' pattern )\* ']'

**patt**[[ [ P1 ] ]] = **patt**[[ P1 :: [ ] ]]

**patt**[[ [ P1 ; P2 ... ] ]] = **patt**[[ P1 :: [ P2 ... ] ]]

**...** : ( ';' pattern )\*



# Tuple patterns

\* pattern ::= pattern ( ',' pattern )+

**patt**[[ P1 , P2 ... ]] =  
    **invert tuple\_prefix**(**patt**[[ P1 ]],  
                          **patt\_tuple**[[ P2 ... ]])

... : ( ',' pattern )\*

**SEMANTICS**    **patt\_tuple**[[ P ... ]] : patt

- **patt\_tuple**[[ P1 ]] =  
    **invert tuple\_prefix**(**patt**[[ P1 ]],  
                          **only**(**tuple\_empty**))

- **patt\_tuple**[[ P1 , P2 ... ]] =  
    **invert tuple\_prefix**(**patt**[[ P1 ]],  
                          **patt\_tuple**[[ P2 ... ]])

tuple < value

# Loose record patterns

```
* pattern ::= '{' field '=' pattern  
           ( ';' field '=' pattern )* '}'
```

```
patt[[ { F1 = P1 ... } ]] =  
  loose_record(patt_map[[ F1 = P1 ... ]])
```

```
... : ( ';' field '=' pattern )*
```

**SEMANTICS**    **patt\_map**[[ F = P ... ]] : map

```
- patt_map[[ F1 = P1 ]] = map1(F1, patt[[ P1 ]])
```

```
- patt_map[[ F1 = P1 ; F2 = P2 ... ]] =  
  map_union(patt_map[[ F1 = P1 ]],  
            patt_map[[ F2 = P2 ... ]])
```



# Declarations

# Pattern matching declarations

**SYNTAX**       $LB : \text{let\_binding}$

**SEMANTICS**     $\text{decl} [[ LB ]] : \text{decl}$

\*  $\text{let\_binding} ::= \text{pattern '=' expr}$

$\text{decl} [[ P = E ]] = \text{match}(\text{expr} [[ E ]], \text{patt} [[ P ]])$

\*  $\text{let\_binding} ::= \text{value\_name pattern+ '=' expr}$

$\text{decl} [[ VN P \dots = E ]] = \text{decl} [[ VN = \text{fun } P \dots \rightarrow E ]]$

$\dots : \text{pattern}^*$



# Top-level phrases

**SYNTAX**       $TP : \text{toplevel\_phrase}$

**SEMANTICS**     $\text{decl}[[ TP ]] : \text{decl}$

\*  $\text{toplevel\_phrase} ::= \text{definition}$

$\text{decl}[[ TP ]] = \text{decl}[[ D ]]$   
when  $TP = D$

\*  $\text{toplevel\_phrase} ::= \text{expr}$

$\text{decl}[[ TP ]] = \text{seq}(\text{print}(\text{expr}[[ E ]]), \text{map\_empty})$   
when  $TP = E$

# Top-level inputs

**SYNTAX**      *TI* : topLevel\_input

**SEMANTICS**    *decl* [[ *TI* ]] : **decl**

\* topLevel\_input ::= topLevel\_phrase\* ';;'

*decl* [[ ;; ]] = **map\_empty**

*decl* [[ *TP* ... ;; ]] =  
    **accum**(*decl* [[ *TP* ]], *decl* [[ ... ;; ]])

... : topLevel\_phrase\*



# Programs

**SYNTAX**       $PR$  : program

**SEMANTICS**     $decl$  [[  $PR$  ]] :  $decl$

\* program ::= toplevel\_input+

$decl$  [[  $TI$  ... ]] =  $scope$ ( $ocaml\_light\_library$ ,  
                                   $decl\_seq$  [[  $TI$  ... ]])

... : toplevel\_input\*

**SEMANTICS**     $decl\_seq$  [[  $TI$  ... ]] :  $decl$

–  $decl\_seq$  [[  $TI$  ]] =  $decl$  [[  $TI$  ]]

–  $decl\_seq$  [[  $TI1$   $TI2$  ... ]] =

$accum$ ( $decl$  [[  $TI1$  ]],  $decl\_seq$  [[  $TI2$  ... ]])

# Funcons

# Computation types

Type **computes** ( $X$ )

Type **comm** = **computes** (**skip**)

Type **decl** = **computes** (**env**)

Type **expr** = **computes** (**value**)

Type **depends** ( $X, Y$ )

Type **abs** = **depends** (**passable**, **value**)

Type **patt** = **depends** (**value**, **env**)

# Value types

<b>atom</b>	< value	<b>passable</b>	= value
<b>boolean</b>	< value	<b>throwable</b>	= value
<b>int</b>	< value	<b>bindable</b>	= value
<b>float</b>	< value	<b>storable</b>	= value
<b>char</b>	< value		
<b>string</b>	< value		
<b>tuple</b>	< value		
<b>record</b>	< value		
<b>variant</b>	< value		
<b>function</b>	< value		
<b>variable</b>	< value		
<b>skip</b>	< value		

```

Funcon abs(patt,expr)           : abs
Funcon accum(env,decl)         : decl
Funcon any : patt
Funcon apply(depends(X,Y),X)    : computes(Y) { val(X), val(Y) }
Funcon apply_function(function,value) : expr
Funcon apply_to_each(proc,list(value)) : comm
Funcon bind(id)                : patt
Funcon bound_value(id)        : computes(bindable)
Funcon catch(expr,abs)        : expr
Funcon close F(abs)           : abs { F(abs) : abs }
Funcon effect(X)              : comm { comp(X) }
Funcon given                  : depends(X,X)
Funcon if_true(boolean,X,X)   : X
Funcon invert F(patt,patt)    : patt { F(X,Y) : Z}
Funcon invert F(patt)        : patt { F(X) : Y}
Funcon invert F               : patt { F : X}
Funcon loose_record(map(field,patt)) : patt
Funcon match(value,patt)      : decl
Funcon only(value)            : patt
Funcon prefer_over(depends(X,Y),depends(X,Y)) : depends(X,Y)
Funcon print(value)           : comm
Funcon recursive(decl)       : decl
Funcon scope(env,X)           : X { comp(X) }
Funcon seq(comm,X)            : X { comp(X) }
Funcon throw(throwable)      : X { comp(X) }
Funcon while_true(computes(boolean),comm) : comm

```

# Procedural abstraction

Funcon **abs**(patt,expr) : abs

Glossary:

`$abs(Patt,Expr)$` abstracts an expression over a given pattern. `$Patt$` is evaluated in the current (dynamic) environment, and `$Expr$` evaluated in the environment generated by `$Patt$`.

Rules:

$$\text{Patt} \rightarrow \text{Patt}'$$

---

**abs**(Patt,Exp)  $\rightarrow$  abs(Patt',Exp)

**abs**(Env,Exp)  $\rightarrow$  **scope**(Env,Exp)

# Binding accumulation

Function **accum**(env, decl) : decl

Glossary:

Elaborates a \$Decl\$ with \$Env\$ overriding the current environment, and then adds elaborations onto that \$Env\$.

Rules:

$$\frac{\text{env map\_over}(\text{Env}, \text{Env1}) \vdash \text{Decl} \rightarrow \text{Decl}'}{\text{env Env1} \vdash \mathbf{accum}(\text{Env}, \text{Decl}) \rightarrow \text{accum}(\text{Env}, \text{Decl}')} \\ \mathbf{accum}(\text{Env}, \text{Env1}) \rightarrow \text{map\_over}(\text{Env1}, \text{Env})$$

# Anonymous pattern

Funcon **any** : patt

Glossary:

The pattern `$any$` matches any value.

Rules:

given Value |- **any** --> map\_empty





# Function application

Funcon **apply\_function**(function, value) : expr

Glossary:

`apply_function(Function, Value)` applies a `Function` to a `Value` and returns the result.

Rules:

**apply\_function**(function(Abs), Value) = **apply**(Abs, Value)

# Execute for each item

Funcon **apply\_to\_each**(proc, list(value)) : comm

Glossary:

`$apply_to_each(Proc,List)$` runs `$Proc$` sequentially on each element of `$List$`.

Rules:

**apply\_to\_each**(Proc, list\_empty) = **skip**

**apply\_to\_each**(Proc, list\_prefix(Value,List)) =  
**seq**(**apply**(Proc, Value), apply\_to\_each(Proc, List))

# Binding pattern

Funcon **bind**(id) : patt

Glossary:

`$bind(Id)$` matches any value, and binds `$Id$` to it.

Rules:

```
given Value |- bind(Id) -->
map_update(map_empty, Id, Value)
```

# Bound value

Funcon **bound\_value**(id) : computes(bindable)

Glossary:

The expression `$bound_value(Id)$` returns the value bound to `$Id$`.

Rules:

$$\frac{\text{map\_select}(\text{Env}, \text{Id}) = \text{Expr}}{\text{env Env} \mid - \text{bound\_value}(\text{Id}) \dashrightarrow \text{Expr}}$$

# Catching thrown values

Funcon **catch**(`expr, abs`) : `expr`

Glossary: The computation `$catch(Expr, Abs)$` evaluates `$Expr$` and returns the result. If `$Expr$` throws an exception, the computation returns `$Abs$` applied to that exception.

Requires: `Throwable < Value`

Rules:

$$\text{Expr} \text{ -- exception none --> Expr'}$$

---

**catch**(`Expr, Abs`) -- exception none --> `catch(Expr', Abs)`

**catch**(`Value, Abs`) --> `Value`

$$\text{Expr} \text{ -- exception some(Throwable) --> Expr'}$$

---

**catch**(`Expr, Abs`) -- exception none -->  
`apply(Abs, Throwable)`

# Abstraction closure

Funcon **close**  $F(\text{abs}) : \text{abs} \{ F(\text{abs}) : \text{abs} \}$

Glossary:

Constructs a statically bound  $F$ .

Rules:

$\text{env Env} \vdash \text{close } F(\text{Abs}) \dashrightarrow F(\text{closure}(\text{Abs}, \text{Env}))$

# Effect

Funcon **effect**(X) : comm { comp(X) }

Glossary:

The computation `$effect(X)$` computes `$X$` until it is a value and disregards the result.

Rules:

$$X \longrightarrow X'$$

---

**effect**(X)  $\longrightarrow$  effect(X')

**effect**(Val)  $\longrightarrow$  **skip**



# Given value

Funcon **given** : depends(X,X)

Glossary:

Returns the given value.

Rules:

$\sim X = \text{null}$

---

given X |- **given** --> X

# Conditional

Funcon **if\_true**(boolean, X, X) : X

Glossary:

At run-time, the computation `if_true(Boolean, X1, X2)` first evaluates `Boolean`. If the result is `true`, it executes `X1`; if the result is `false`, it executes `X2`.

Rules:

**if\_true**(true, X1, X2) = X1

**if\_true**(false, X1, X2) = X2

# Deconstruction

```
Funcon invert F(Patt,Patt) : Patt { F(X,Y) : Z}
Funcon invert F(Patt)      : Patt { F(X) : Y}
Funcon invert F           : Patt { F : X}
```

Glossary: The pattern `invert F(Patt1,Patt2)` matches `F(X,Y)` when `X` matches `Patt1` and `Y` matches `Patt2`. We also include unary and nullary versions.

Rules:

```
              given X |- Patt1 --> Patt1'
-----
given F(X,Y) |- invert F(Patt1,Patt2) -->
invert F(Patt1',Patt2)

              given Y |- Patt2 --> Patt2'
-----
given F(X,Y) |- invert F(Patt1,Patt2) -->
invert F(Patt1,Patt2')
```

```
given F(X,Y) |- invert F(Env1,Env2) --> map_over(Env1,Env2)

              ~ Value = F(X,Y)
-----
given Value |- invert F(Patt1,Patt2) -- not_in_domain true --> stuck

              given X |- Patt --> Patt'
-----
given F(X) |- invert F(Patt) --> invert F(Patt')
```

```
given F(X) |- invert F(Env) --> Env

              ~ Value = F(X)
-----
given Value |- invert F(Patt) -- not_in_domain true --> stuck

given F |- invert F --> map_empty

              ~ Value = F
-----
given Value |- invert F -- not_in_domain true --> stuck
```

# Loose record pattern

Funcon **loose\_record**(map(field,patt)) : patt

## Glossary:

`$loose_record(Map,Pattern)$` matches any `$Record$` such that whenever `$map_select(Map,Field) = Patt$`, the `$Field$` projection of `$Record$` matches `$Patt$`.

## Rules:

```
given Record |- loose_record(map_empty) --> map_empty
```

```
given Record |- loose_record(map_prefix(Field,Patt,Map1)) -->  
accum(match(Record,record_item(Field,Patt)),match(Record,loose_record(Map1)  
))
```

~ Value : Record

---

```
given Value |- loose_record(Map) -- not_in_domain true --> stuck
```

# Pattern matching declaration

Funcon **match**(value,patt) : Decl

Glossary:

`$match(Value, Patt)` tries to match `$Value` against `$Patt`, computing the resulting bindings.

Rules:

$$\text{given Value } |- \text{ Pat } \longrightarrow \text{ Pat '}$$

---

$$\text{given Value1 } |- \text{ match(Value, Pat) } \longrightarrow \text{ match(Value, Pat ')}$$
$$\text{match(Value, Env) } \longrightarrow \text{ Env}$$

# Constant pattern

Funcon **only**(value) : patt

Glossary:

`$only(Groundval)$` matches only the data value `$Value$`.

Rules:

Value : groundval

---

given Value |- **only**(Value) --> map\_empty

Value1 : groundval, Value2 : groundval,  
~ Value1 = Value2

---

given Value1 |- **only**(Value2) -- not\_in\_domain true -->  
stuck

# Abstraction preference

Funcon **prefer\_over**(depends(X,Y), depends(X,Y)) : depends(X,Y)

Glossary:

`prefer_over(Dep1,Dep2)` evaluates `Dep1` on the `given` argument. If it's undefined, `Dep2` is applied to the `given` argument.

Rules:

$$\text{Dep} \text{ -- not\_in\_domain false --> Dep'}$$

---

**prefer\_over**(Dep,Dep1) -- not\_in\_domain false --> prefer\_over(Dep',Dep1)

**prefer\_over**(Val,Dep) --> Val

$$\text{Dep} \text{ -- not\_in\_domain true --> Dep'}$$

---

prefer\_over(Dep,Dep1) -- not\_in\_domain false --> Dep1

# Print

Funcon **print**(value) : comm

Glossary:

The computation `$print(Value)$` prints `$Value$` to output.

Rules:

**print**(Value) -- output `list_prefix(Value, list_empty)` -->  
skip



# Recursive declaration

Funcon **recursive**(decl) : decl

Glossary:

`$recursive(Env)$` evaluates computation entries in `$Env$` with respect to itself, and returns the resulting environment of values. 1. Compute forwards and get resulting map. 2. Duplicate map. 3. From one copy: construct environment of forwards. 4. Evaluate entries of `$Env$` within that environment of forwards. 5. From second copy: for each id, connect that forward to the current bound value. NB: Declaration is evaluated twice in current model (once to find out which identifiers are bound, again to evaluate in the environment with those identifiers bound to forwards.)

Rules:

**recursive**(Decl) = reclose(generate\_forwards(Decl),Decl)

Funcon generate\_forwards(env) : computes(map(id,fwd))

Rules:

generate\_forwards(map\_empty) = map\_empty

next\_forward(Fo) = Fwd

-----  
(generate\_forwards(map\_prefix(Id,Expr,Env)), forwards Fo) --> (map\_prefix(Id,Fwd,generate\_forwards(Env)), forwards map\_update(Fo,Fwd,undefined))

Funcon reclose(map(id,fwd),decl) : decl

Rules:

fwd\_env(MapIdFwd) = Env1

-----  
reclose(MapIdFwd,Decl) --> accum(Env1,accum(Decl,seq(tie(MapIdFwd),Env1)))

Operation fwd\_env(map(id,fwd)) : env

Rules:

fwd\_env(map\_empty) = map\_empty

fwd\_env(map\_prefix(Id,Fwd,Map)) = map\_prefix(Id,forward(Fwd),fwd\_env(Map))

Funcon tie(map(id,fwd)) : cmd

Rules:

tie(map\_empty) = skip

env Env |- (tie(map\_prefix(Id,Fwd,Map)), forwards Fo) --> (tie(Map), forwards map\_update(Fo,Fwd,map\_select(Env,Id)))

# Scope restriction

Funcon **scope**(env, X): X { comp(X) }

Glossary:

The computation `$scope(Env, X)$` runs `$X$` with respect to the bindings in `$Env$`.

Rules:

$$\frac{\text{env map\_over}(\text{Env}, \text{Env1}) \mid\text{- } X \text{ --> } X'}{\text{env Env1} \mid\text{- } \mathbf{scope}(\text{Env}, X) \text{ --> } \text{scope}(\text{Env}, X')}$$
$$\mathbf{scope}(\text{Env}, \text{Val}) \text{ --> Val}$$

# Sequential execution

Function **seq**(comm, X) : X { comp(X) }

Glossary:

The computation `seq(Comm, X)` first evaluates `Comm`, and then `X` is evaluated to a result.

Rules:

$$\frac{\text{Comm} \longrightarrow \text{Comm}'}{\text{seq}(\text{Comm}, X) \longrightarrow \text{seq}(\text{Comm}', X)}$$
$$\text{seq}(\text{skip}, X) \longrightarrow X$$

# Throw a value

Funcon **throw**(throwable) : X { comp(X) }

Uses: **stuck**

Glossary:

The computation `$throw(Throwable)$` throws the exception `$Throwable$`.

Rules:

**throw**(Throwable) -- exception some(Throwable) --> **stuck**

# Loop

Funcon **while\_true**(computes(boolean), comm) : comm

Uses: if\_true, seq, skip

Glossary:

The computation `$while_true(Expr, Comm)$` repeatedly runs `$Comm$` as long as the condition `$Expr$` holds.

Rules:

**while\_true**(Expr, Comm)  $\rightarrow$   
**if\_true**(Expr, **seq**(Comm, **while\_true**(Expr, Comm)), **skip**)

# Ad hoc abbreviation

Funcon **ocaml\_light\_library** : env

Rules:

```
ocaml_light_library --> { id('+') |--> function(curry(abs int_plus)),
  id('-') |--> function(curry(abs int_minus)),
  id('*') |--> function(curry(abs int_times)),
  id('/') |--> function(curry(abs int_divide)),
  id('mod') |--> function(curry(abs int_modulo)),
  id('land') |--> function(curry(abs int_and)),
  id('lor') |--> function(curry(abs int_or)),
  id('lxor') |--> function(curry(abs int_xor)),
  id('lsl') |--> function(curry(abs int_lshift)),
  id('lsr') |--> function(curry(abs int_rshift)),
  id('asr') |--> function(curry(abs int_arshift)),
  id('+.') |--> function(curry(abs float_plus)),
  id('-.') |--> function(curry(abs float_minus)),
  id(*.') |--> function(curry(abs float_times)),
  id('/.') |--> function(curry(abs float_divide)),
  id('**') |--> function(curry(abs float_exp)),
  id('@') |--> ol_append,
  id('^') |--> function(curry(abs string_append)),
  id('!!') |--> function(abs deref),
  id(':=') |--> function(curry(abs assign)),
  id('=') |--> function(curry(abs equal)),
  id('<>') |--> function(curry(compose(abs not,abs equal))),
  id('==') |--> function(curry(abs equal)),
  id('!=') |--> function(curry(compose(abs not,abs equal))),
  id('<') |--> function(curry(abs less)),
  id('>') |--> function(curry(abs greater)),
  id('<=') |--> function(curry(abs less_equal)),
  id('>=') |--> function(curry(abs greater_equal)),
  id('not') |--> function(abs not),
  id('sin') |--> function(abs float_sin),
  id('pi') |--> float_pi,
  id('float') |--> function(abs int_to_float) }
```

# *Preliminary* tool support

IDE :

- ▶ **Spoofax**

Parsing, translation to funcons :

- ▶ **ASF+SDF**

Funcon interpretation :

- ▶ **Prolog**

# PLANCOMPS

Programming Languages  
(incl. DSLs)

C#

Java

...



Components and their Specifications



fundamental constructs  
'funcons'

reusable